#### Understanding Multi-Key Security Degradation

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<sup>1</sup>ESAT/COSIC, KU Leuven and iMinds, Belgium

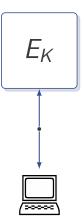
<sup>2</sup>Computer Science Department, UC Davis, USA

<sup>3</sup>Information Security Group, Royal Holloway, University of London, UK

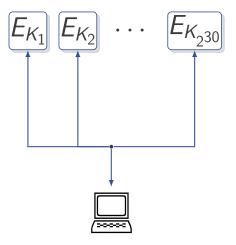
September 30, 2016

Single-Key vs. Multi-Key

1. Single-key setting usually analyzed

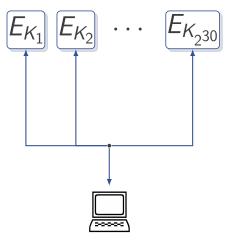


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- 1. Single-key setting usually analyzed
- 2. Multi-key setting is practically important
- Example: AES-GCM used in TLS, hundreds of millions of keys used

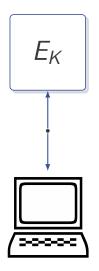
Folklore Result and Multi-Key Degradation

# Multi-Key Success Probability

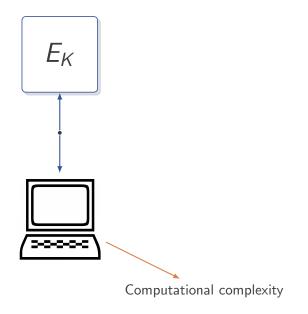


# $u \times$ Single-Key Success Probability

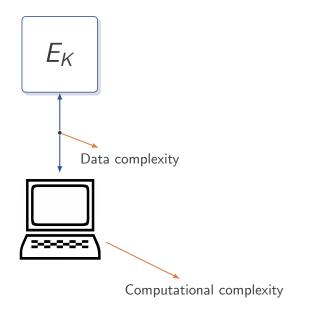
## Example : Block Ciphers



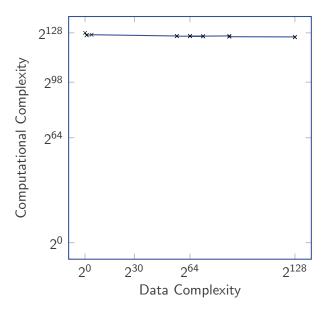
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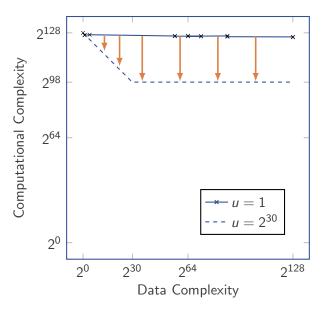
## Example : Block Ciphers



Example: AES128 Key Recovery



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Precomputation:

 $(L, E_L(P))$ 

Precomputation:

Queries:

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 $E_{K_1}(P), E_{K_2}(P), \ldots, E_{K_u}(P).$ 

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Biham 2002 Information Processing Letters

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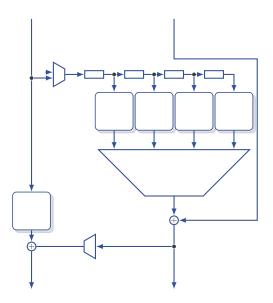
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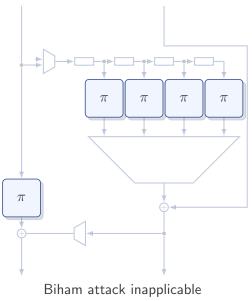
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Biryukov, Mukhopadhyay, Sarkar, SAC 2005: time-memory-data trade-off

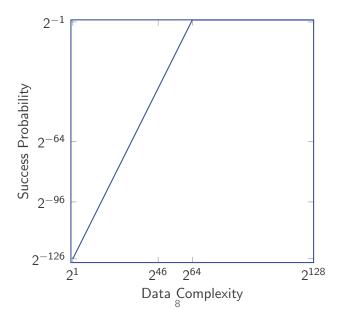
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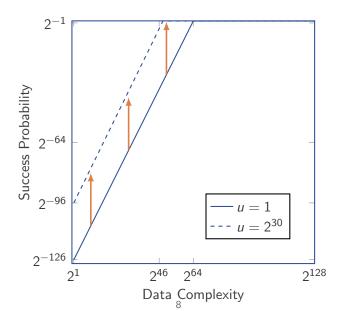
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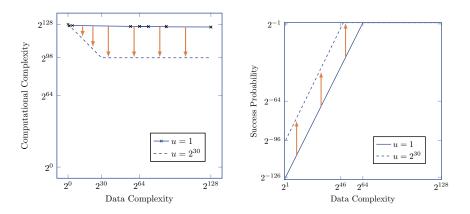
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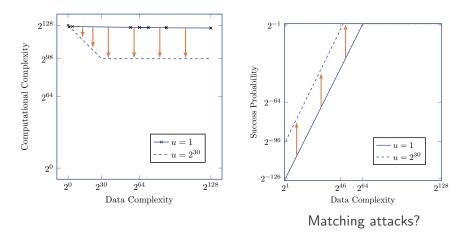
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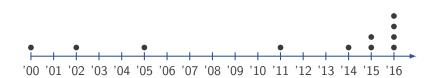


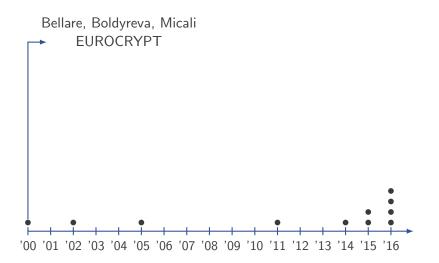
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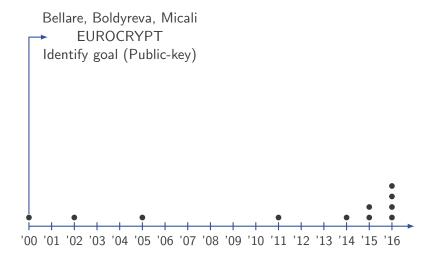


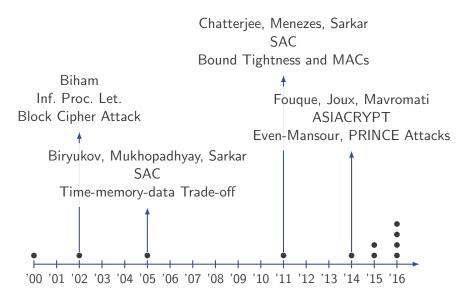
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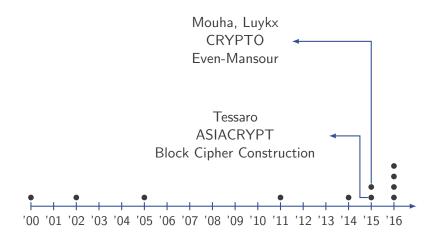


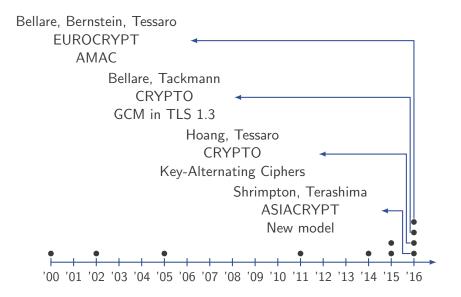


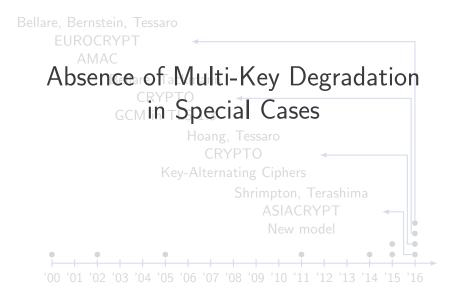


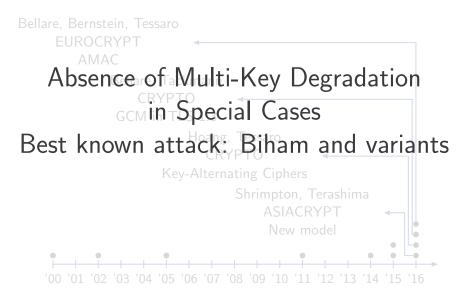






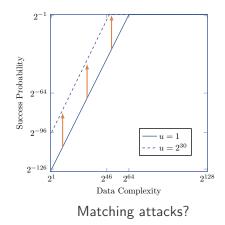






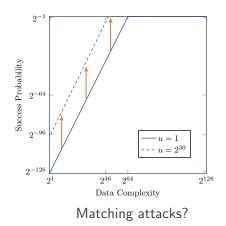
#### Our Work

#### 1. Set out to understand gap



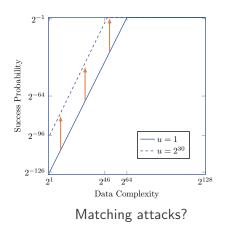
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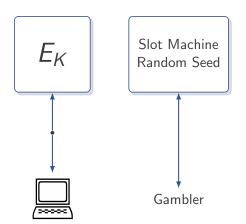
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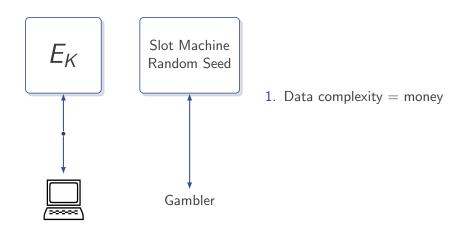


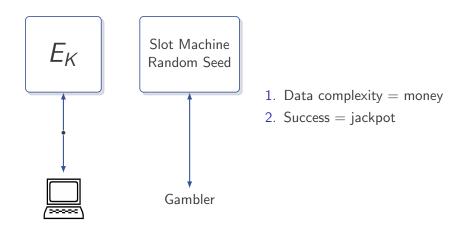
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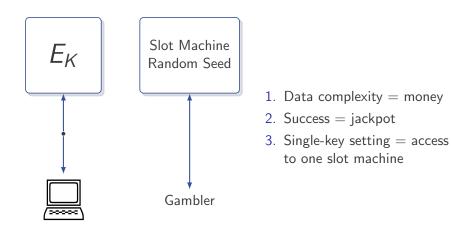
- 1. Set out to understand gap
- Characterization of multi-key setting: necessary and sufficient condition for degradation
- 3. Proved in abstract setting, applied to GCM











## Slot Machine Scenario



500 coin budget

500 coin budget, 100 slot machines

# 500 coin budget, 100 slot machines all 500 coins on 1 machine

VS

500 coins distributed somehow over 100 machines

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"spending 500 coins on 100 slot machines gives you factor 100 higher success than spending 500 coins on one slot machine"

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Counter-intuitive, ever possible? Yes.

- 1. Assume slot machine is "lucky" with some probability
- 2. One slot machine: either lucky or not.
- 3. One hundred slot machines: find lucky machine, focus on that one

Weak keys

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Weak-key recovery: computational complexity 2<sup>16</sup>, data complexity 2.

Table: Midori64 key recovery

	u = 1	$u = 2^{16}$
Computational cost	2 <sup>16</sup>	2 <sup>17</sup>
Data cost	2	<i>u</i> + 2
Success Estimate	$2^{-96}$	$2^{-80}$

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- 2. Exploit Weak instances if present
- 3. What else can happen?

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# Setting 1: possibility of interacting with lucky machine $\Rightarrow$ jackpot probability might be higher

#### Translation to Oracles and Games

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Setting 1 Adversarial Success  $\leq$  Setting 2 Adversarial Success given transcript is satisfied

for all transcripts below some cost

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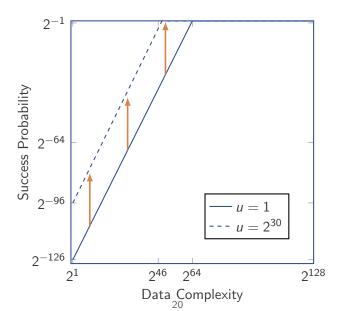
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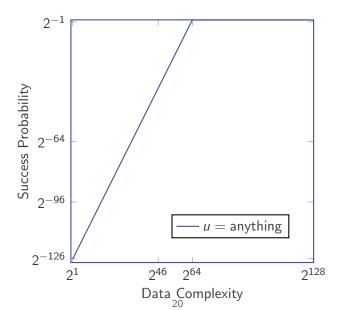
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This is only intuition, proper formalization introduces subtleties! (Information-theoretic setting, adversaries must be optimal, queries are bounded,...)

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Thank you for your attention.